

7th Central and Eastern European
Software Engineering Conference
in Russia - CEE-SECR 2011

October 31 – November 3, Moscow



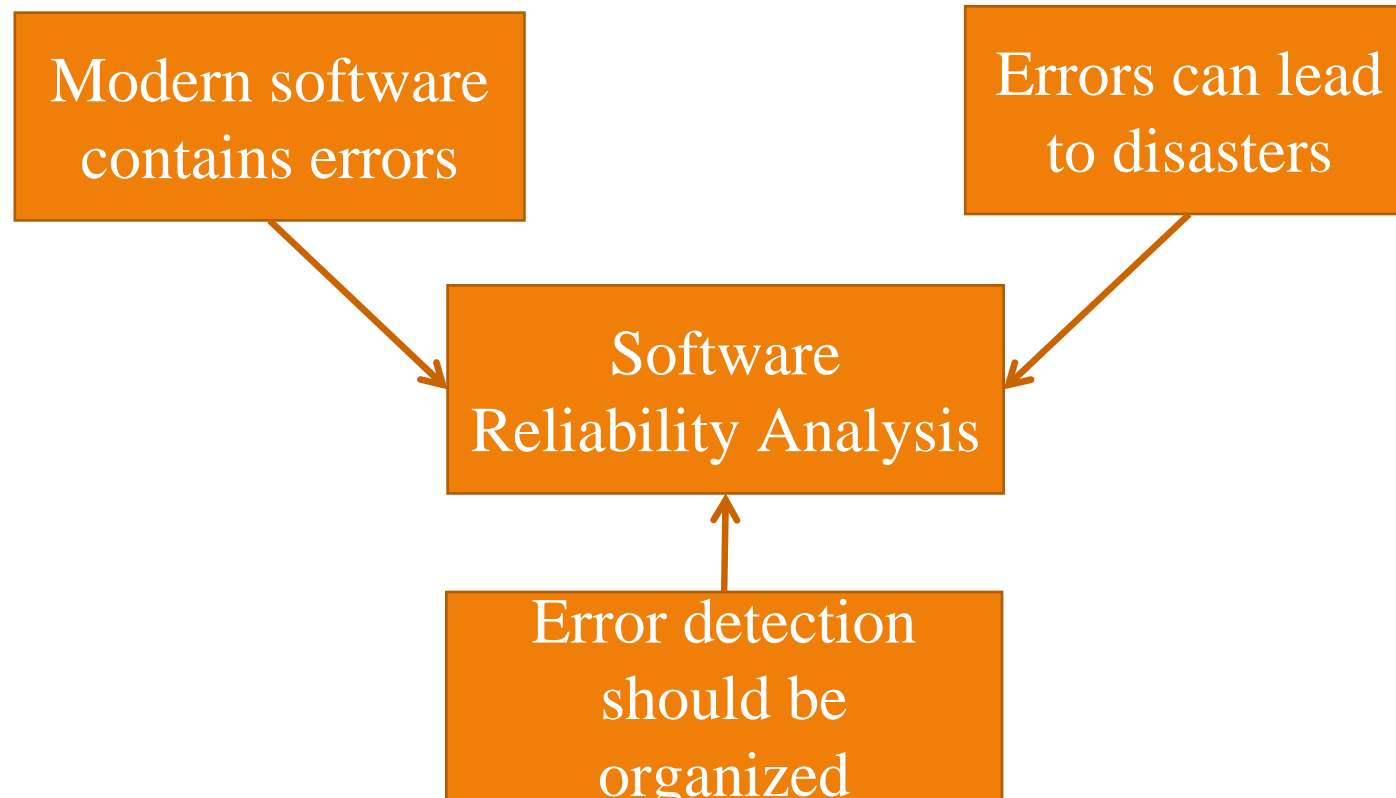
Software Reliability Estimation Based on Static Error Detection

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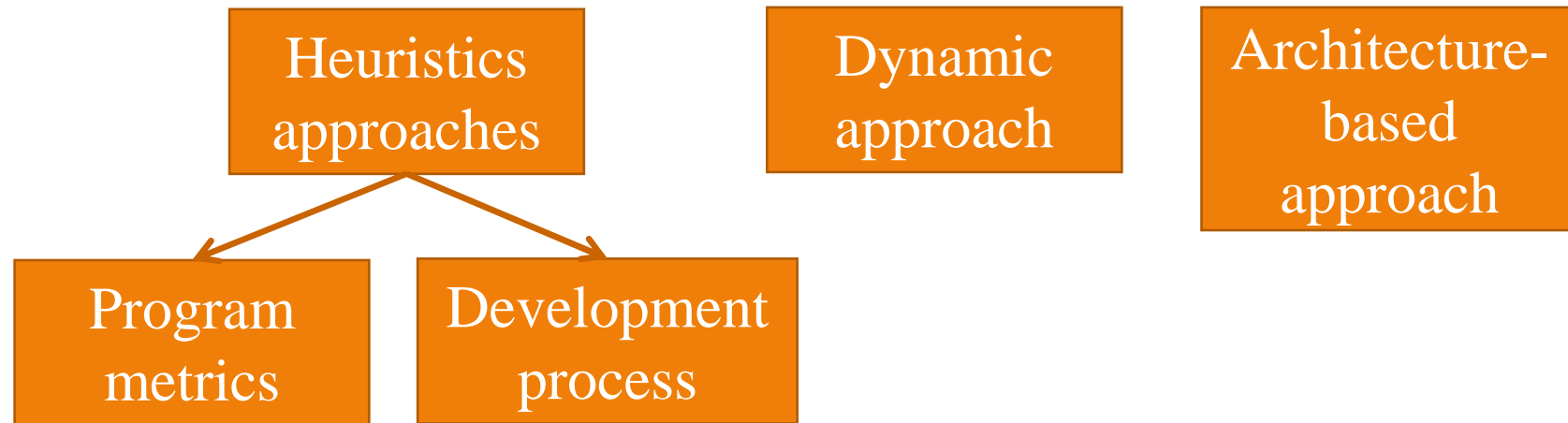


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Importance of Software Reliability Analysis



Known Approaches



Known Approaches – Program Metrics

- Based on simple code properties, such as
 - number of statements
 - number of conditions
 - number of loops
 - number of functions
 - ...

Known Approaches – Development Process Metrics

- Based on development process properties, such as
 - duration of development
 - number & qualification of developers
 - number & qualification of testers
 - methodology used
 - automation tools used

Known Approaches – Others

- Runtime
 - Based on failures observed at run-time
- Architecture-based
 - Based on known reliability of program components

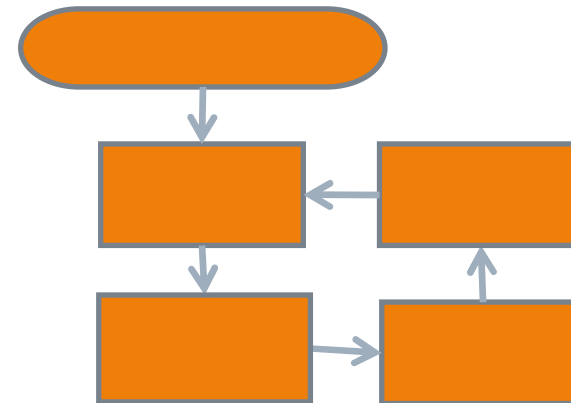
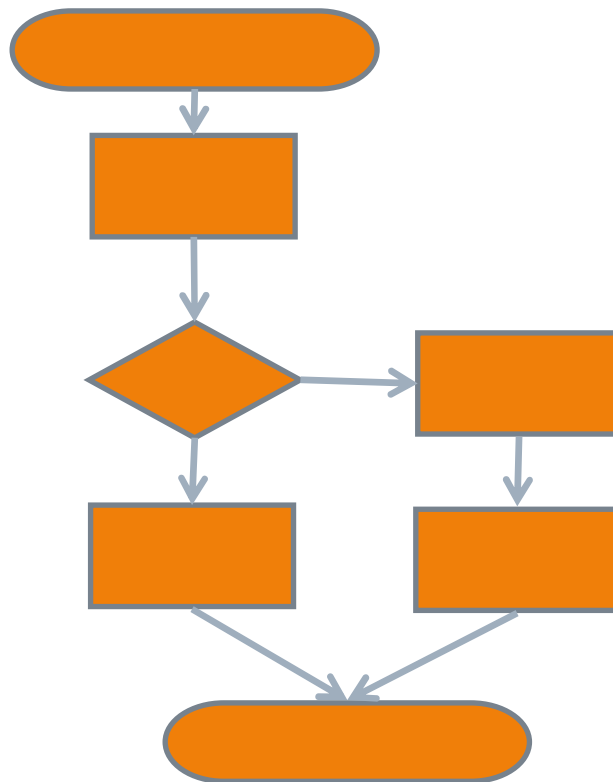
Our Approach

- Based on source code static analysis
- Delivers
 - Ranking of errors (based on failure probability)
 - Reliability characteristics
- Limitations
 - Single-threaded C programs
- Error types
 - uninitialized variable use
 - incorrect pointer dereference
 - pointer out of bounds

Features of Our Approach

- Analysis of a program model
- Analysis of all possible execution paths
- Advantages
 - Reliability estimations is based on real errors
 - Results are applicable for any exploitation conditions
 - Makes debugging more effective
- Drawbacks
 - Does not consider quantitative time
 - Does not consider normal program exploitation
 - Execution path probability estimation
 - False positives problem

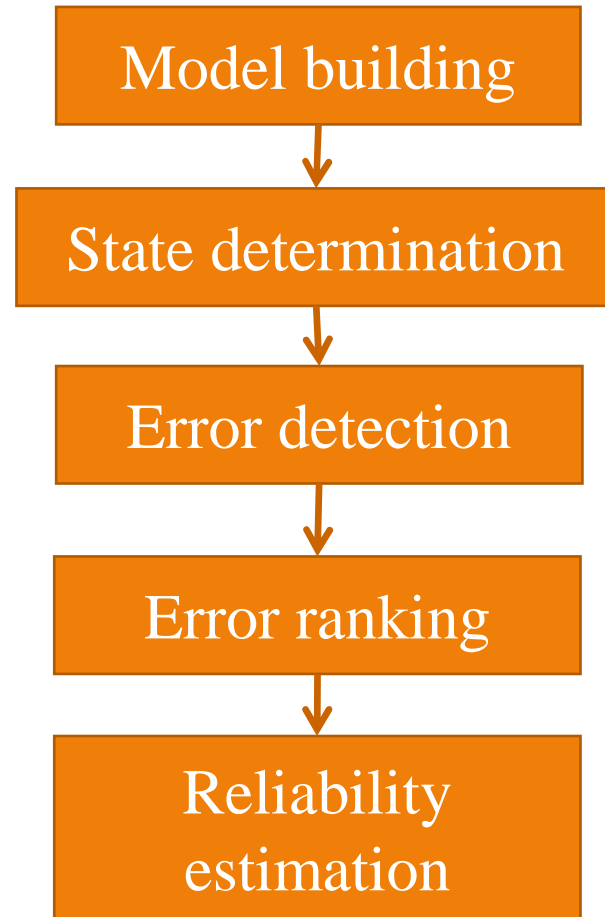
Program Classes



Reliability characteristics used

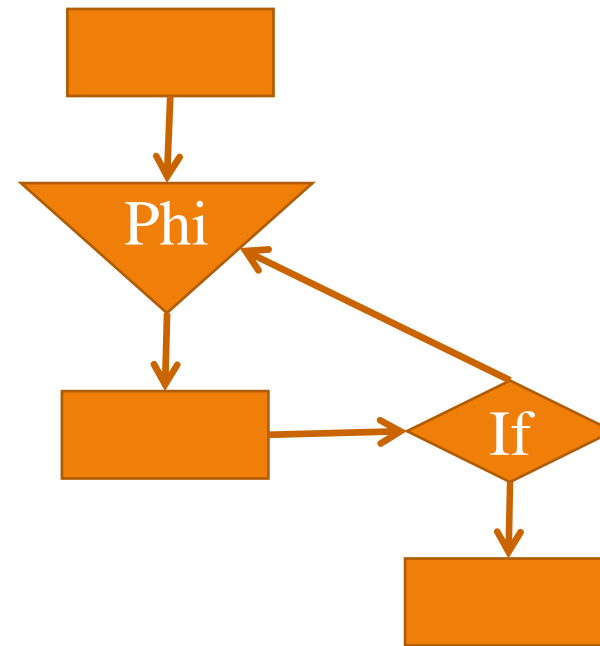
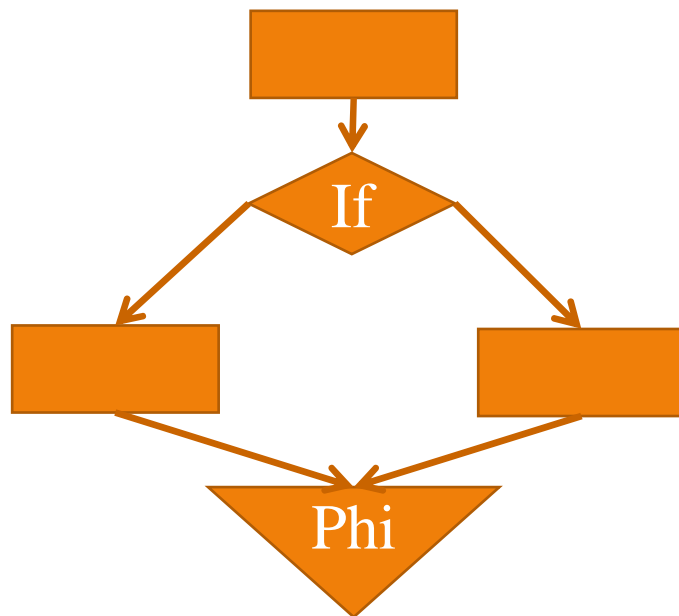
- Computational programs
 - Probability of whole program successful execution $P(\infty)$
- Server programs
 - Probability of n statements successful execution $P(n)$
 - Mean executed statement number before failure \bar{n}

Algorithms



Program Model Features

- Control flow graph
- Three-operand assignment form $A = B \text{ op } C$
- If and Phi statements



State Determination Algorithms

- State representation
- Control flow analysis
 - Statement analysis
 - Sequential
 - If statement analysis
 - Phi statement analysis
 - Loop analysis
 - Interprocedural analysis

Program State Representation

- Based on objects, values, and probabilities
 - set of triples $Q = \{(o_j, v_k, p_{jk})\}$
 - state probability $P(Q)$
- Object values
 - intervals
 - pointers
 - resource descriptors

Probability normalization

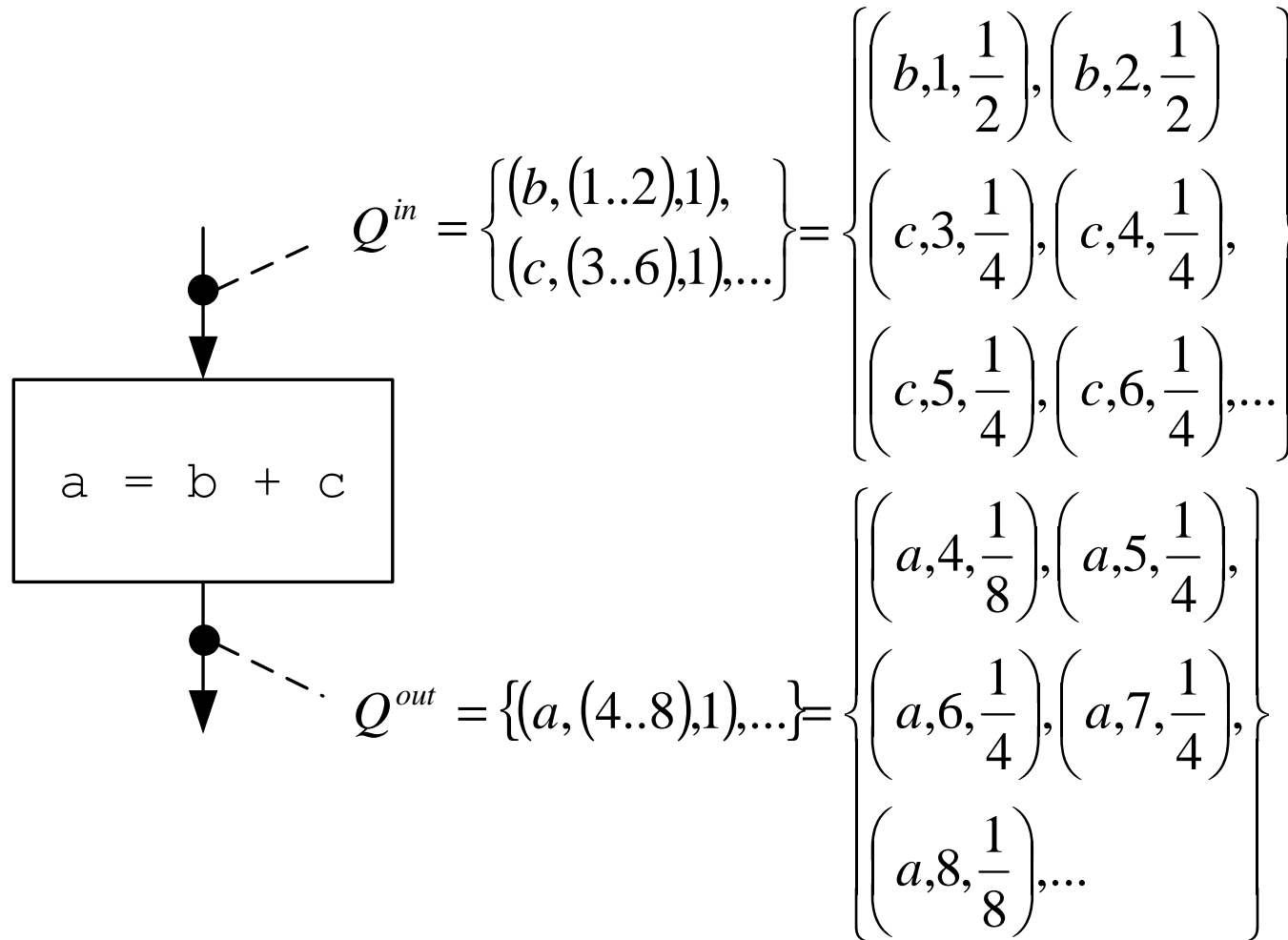
- Control flow normalization

$$\sum_{\forall Q_j^{in} \in \text{Input}(s)} P(Q_j^{in}) = \sum_{Q_j^{out} \in \text{Output}(s)} P(Q_j^{out})$$

- State normalization

$$\forall o_j : \exists (o_j, v_k, p_{jk}) \in Q \Rightarrow \sum_{\forall (o_j, v_k, p_{jk}) \in Q} p_{jk} = P(Q)$$

Sequential Statement Analysis



If Statement Analysis

- True and false combination consideration

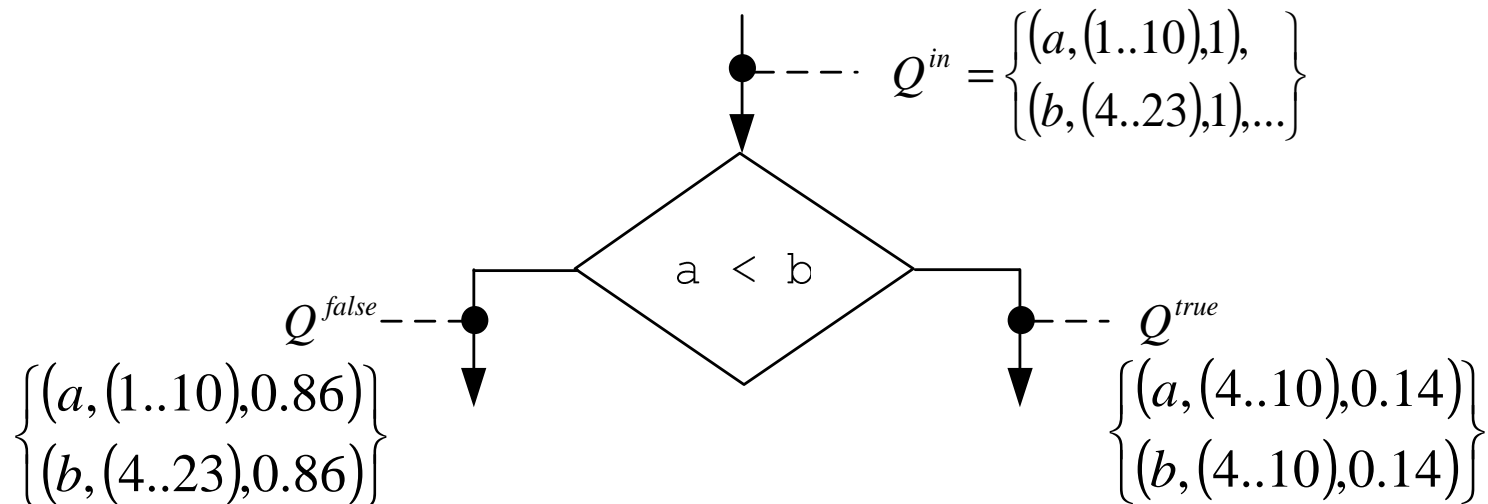
$$P(Q^{true}) = \sum_{c \in C^{true}} \prod_{(o_j, v_k, p_{jk}) \in c} p_{jk},$$

$$P(Q^{false}) = \sum_{c \in C^{false}} \prod_{(o_j, v_k, p_{jk}) \in c} p_{jk}.$$

- Normalization of state probabilities
- Normalization of non-affected triples probabilities

If Statement Analysis Example

- 172 combinations where $a < b$
- 28 combinations where $a \geq b$



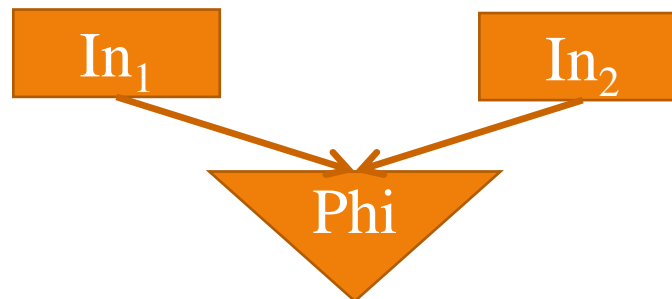
- Normalization: 0.86 for true, 0.14 for false

Phi Statement Analysis

- Identical triples are added together

$$\forall o_j, v_k : (o_j, v_k, p_{jk}) \in Q_1^{in}, (o_j, v_k, r_{jk}) \in Q_2^{in} \Rightarrow (o_j, v_k, p_{jk} + r_{jk}) \in Q^{out}$$

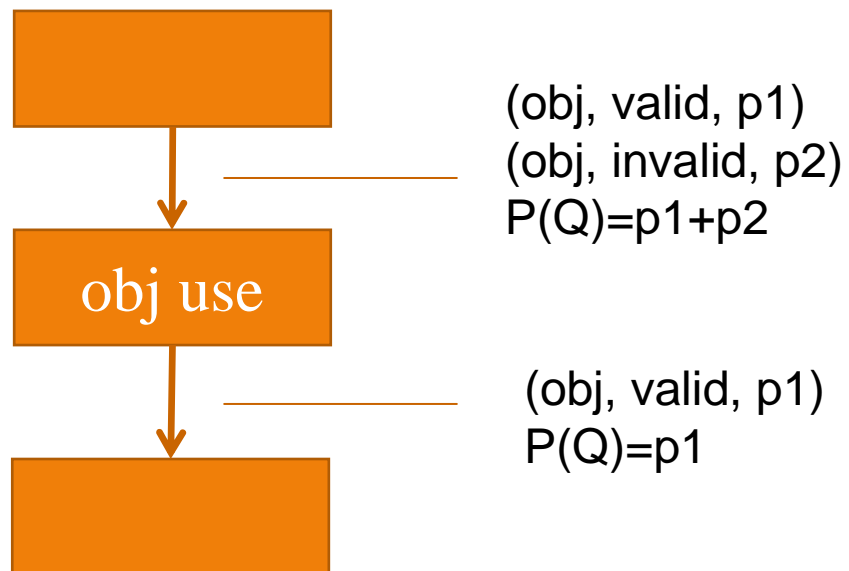
- Control flow normalization $P(Q^{out}) = P(Q_1^{in}) + P(Q_2^{in})$



Error Detection

- Based on incorrect values in state
 - uninitialized variable use $(o_j, v_{noninit}, p_k)$
 - pointer dereference $(o_j, v_{noninit}, p_k)$
 (o_j, v_{null}, p_k) $(o_j, v_{invalid}, p_k)$
 - out of bounds $(o_i, (o_j, offset_j), p_k)$
 - correct if $0 \leq offset_j < sizeof(o_j)$
 - otherwise error is detected

Error Inhibition



Error Ranking

- Errors are sorted according to probability of occurrence
 - Most dangerous errors can be corrected first
- Probabilities are summarized for same errors in the same statement

Overall reliability estimation

- probability of successful execution

$$P(n) = \sum_{\text{end statements}} P(Q)$$

- probability of n statements successful execution

$$P(n) = \sum_{n \text{ statements executed}} P(Q)$$

- mean executed statements number before failure

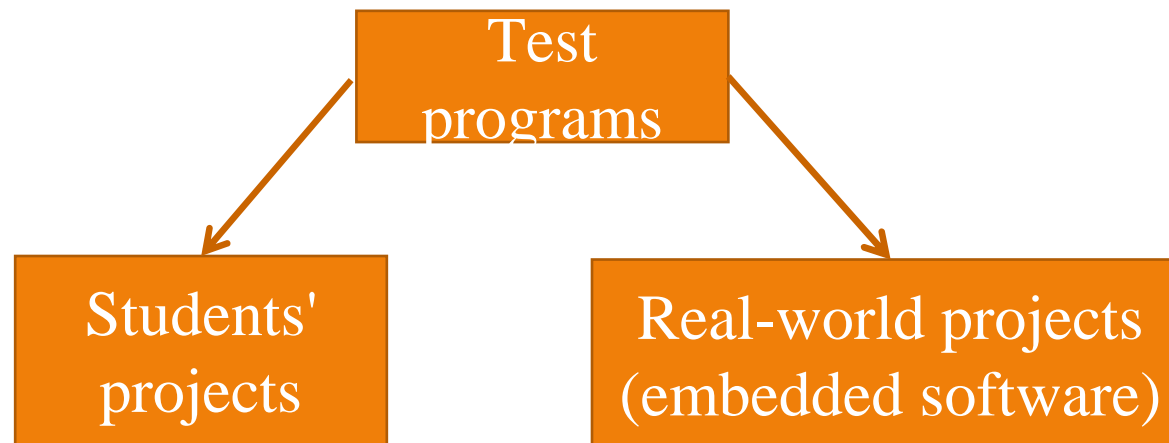
$$\bar{n} = \sum_{n=0}^{n_{\max}} (P(n) - P(n+1)) \cdot n$$

Implementation

- AEGIS static analyzer
 - analysis of C/C++ source code
 - interval, points to, resource analysis
 - loop & interprocedural analysis
 - spread range of program errors detected
- Results
 - error ranking table
 - $P(n)$ table
 - $P(\infty)$
 - mean executed statements number before failure

Experiments made

- Purpose
 - Testing of our approach
 - Debugging example



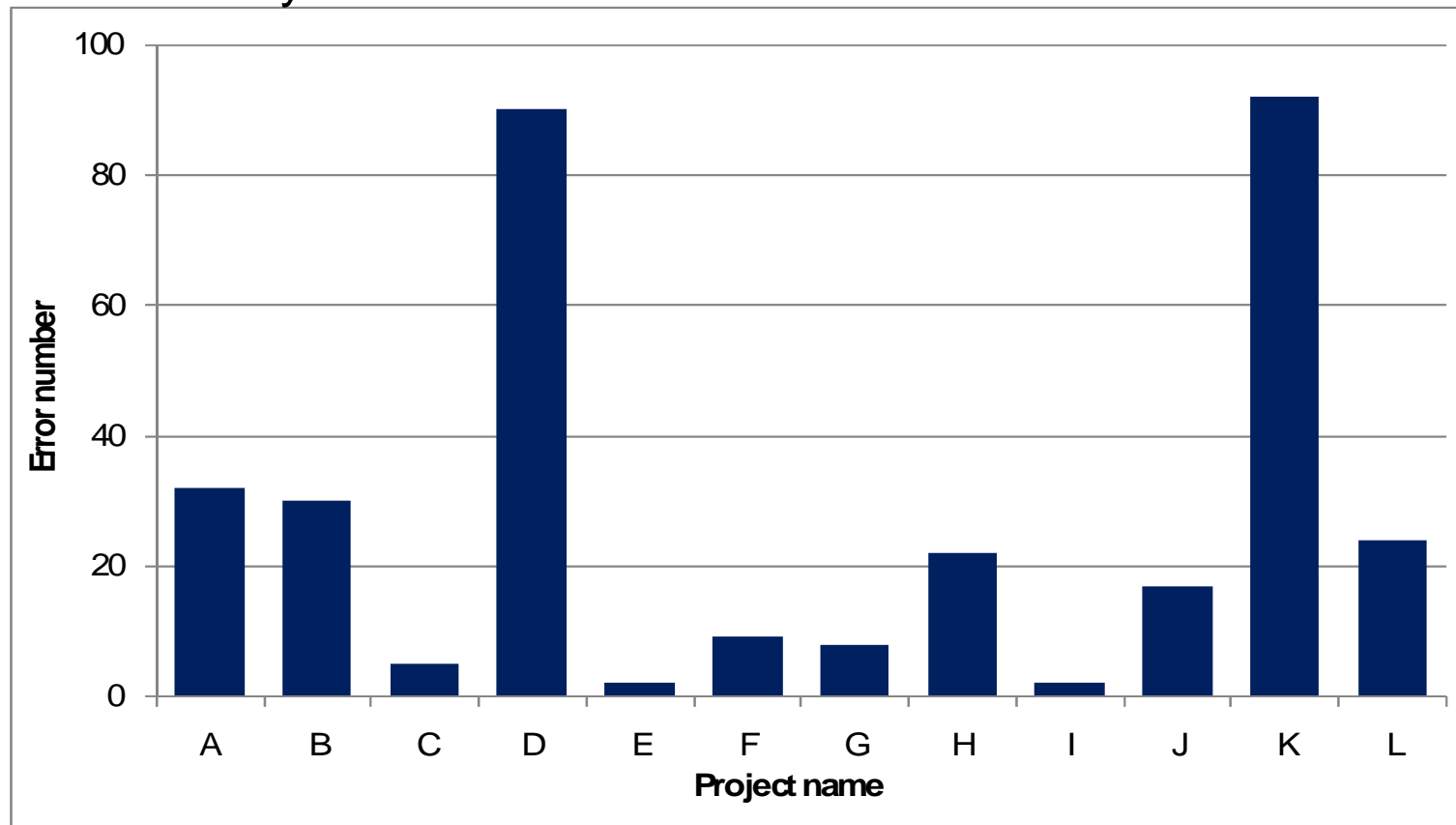
Sample of reliability analysis

```
while (!(feof(f))) // 0.5
{
    i = t = 0;
    // Failure in one of three cases
    prov(&t, strlen(st), st);
}
```

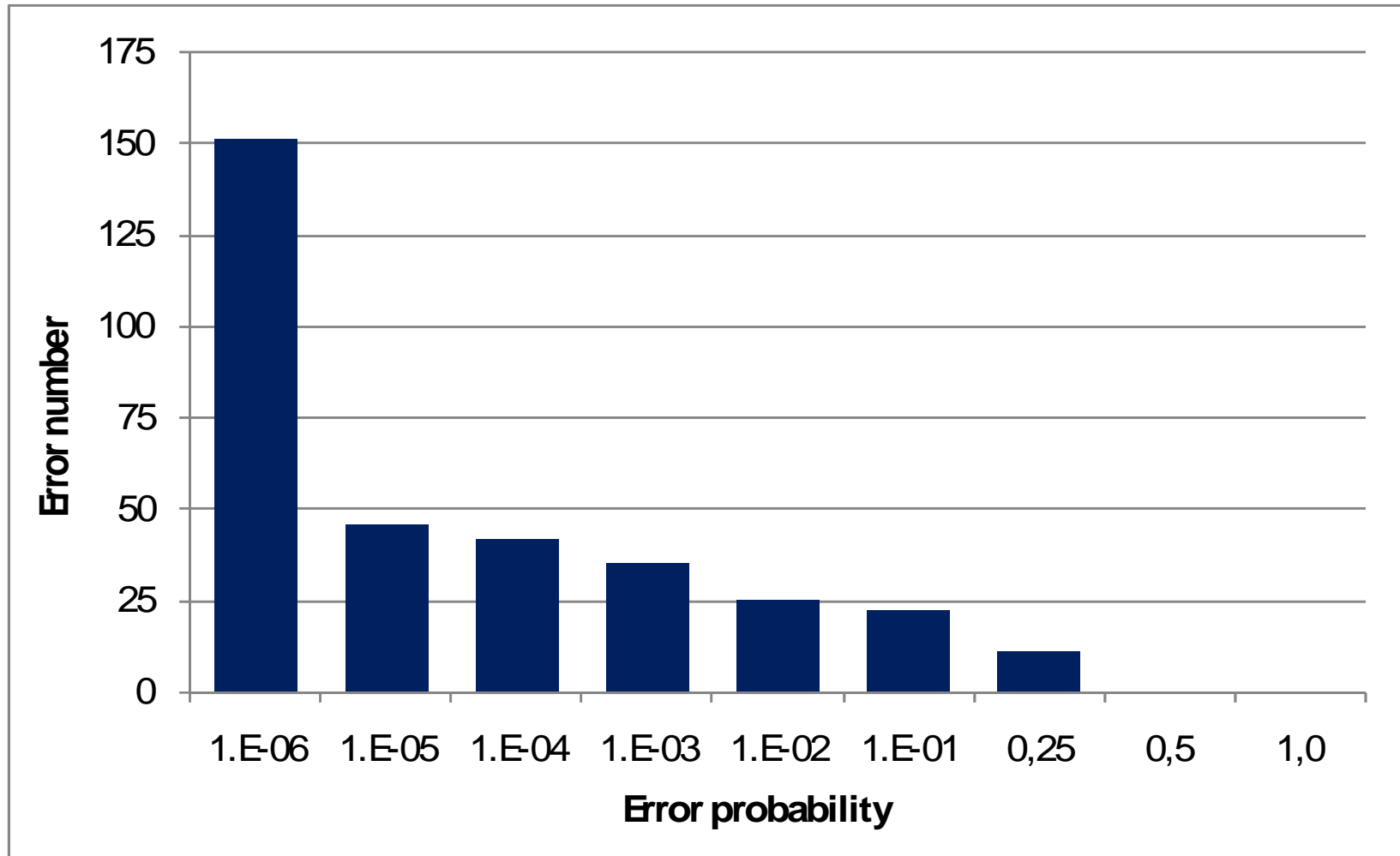
- Probability of successful execution is
 $0.75 = 0.5 + 0.5 * 0.33 + 0.5 * 0.33^2 + \dots$

Amount of errors in real-world projects

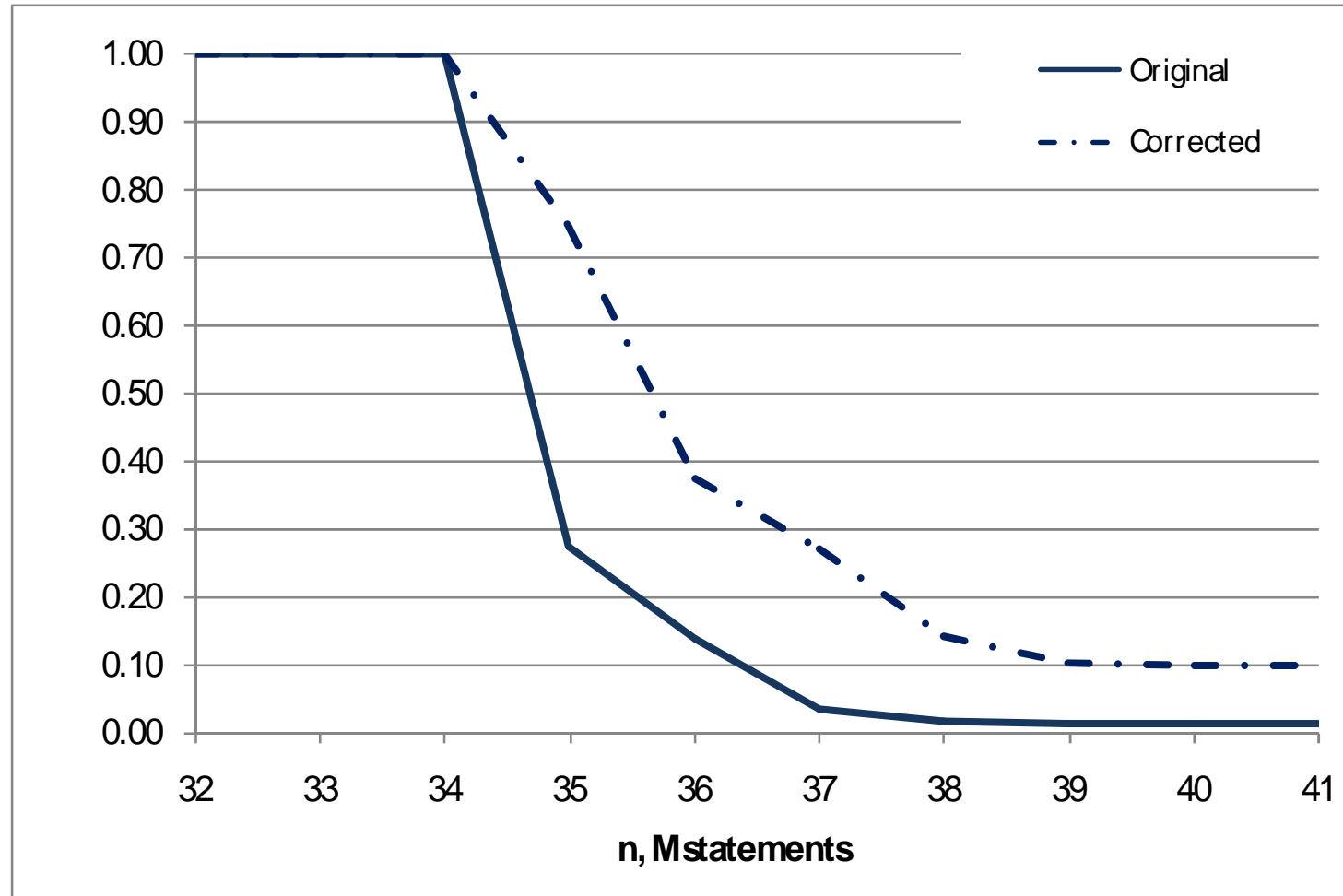
- More than 500 errors, 2/3 of considered types
- Density about 0.8/1KLOC



Distribution of error number



Debugging results



Directions for Future Work

- Reliability estimation
 - Annotations for path probability estimations
 - Run-time analysis for path probability estimation
 - Execution time estimation
- Static analysis itself
 - Soundness & precision
 - Parallel program analysis
 - Annotations for functional error detection

Conclusion

- Approach for software reliability estimation
 - based on error detection using static analysis
- Implementation in AEGIS tool (prototype)
 - ranking of errors by the probability of occurrence
 - probability of successful execution
 - probability of N statement successful execution
 - mean number of executed statements before failure

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